

Online Supplement to “Resource Allocation via Message Passing”

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A. Proofs of Existence Theorems

In the following, a function $f: \mathcal{S} \rightarrow \mathbb{R}$ with domain $\mathcal{S} \subset \mathbb{R}^n$ is said to be Lipschitz with constant L if

$$|f(x) - f(y)| \leq L\|x - y\|, \quad \forall x, y \in \mathcal{S}.$$

Theorem 1. *Assume that the utility functions are Lipschitz continuous. Then, a message-passing equilibrium exists.*

Proof. Let L be a Lipschitz constant that applies to all utility functions. Suppose each message in the set V is Lipschitz continuous with Lipschitz constant L . Consider the message from an activity a to a resource $r \in \mathcal{R}(a)$. Define $\mathcal{X}^{a \setminus r} \triangleq \prod_{r' \in \mathcal{R}(a) \setminus r} \mathcal{X}_r$ to be the space of consumption bundles for activity a , excluding resource r . Without loss of generality, assume that $(FV)_{a \rightarrow r}(x'_{ar}) \geq (FV)_{a \rightarrow r}(x_{ar})$. Then, for some $z' \in \mathcal{X}^{a \setminus r}$,

$$\begin{aligned} (FV)_{a \rightarrow r}(x'_{ar}) - (FV)_{a \rightarrow r}(x_{ar}) &= u_a(x'_{ar}, z') + \sum_{r' \in \mathcal{R}(a) \setminus r} V_{r' \rightarrow a}(z'_{ar'}) \\ &\quad - \max_{z \in \mathcal{X}^{a \setminus r}} \left(u_a(x_{ar}, z) + \sum_{r' \in \mathcal{R}(a) \setminus r} V_{r' \rightarrow a}(z_{ar'}) \right) \\ &\leq u_a(x'_{ar}, z') - u_a(x_{ar}, z') \leq L|x'_{ar} - x_{ar}|. \end{aligned}$$

Hence, the message $(FV)_{a \rightarrow r}(\cdot)$ is Lipschitz continuous with Lipschitz constant L . A similar proof applies to $(FV)_{r \rightarrow a}(\cdot)$.

Let \mathcal{S} be the collection of message sets V for which each message equals zero at zero and is Lipschitz continuous with Lipschitz constant L . Note that \mathcal{S} is convex, closed, and bounded (under the supremum norm). \mathcal{S} is subset of the set of continuous functions from a compact, finite dimensional metric space to itself. Hence, \mathcal{S} is compact under the supremum norm by

the Arzelà-Ascoli theorem. The operator H maps \mathcal{S} to \mathcal{S} continuously with respect to the supremum norm. It follows from the Schauder fixed point theorem that a message-passing equilibrium exists. \square

Theorem 3. *There exists a message-passing equilibrium with concave and Lipschitz continuous messages.*

Proof. The proof follows by a modification of the proof of Theorem 1: define the set \mathcal{S}' to be the collection of message sets $V \in \mathcal{S}$ which are also concave. Since the operator H involves maximization of a concave function over a convex set, if $V \in \mathcal{S}'$, then HV is also concave hence $HV \in \mathcal{S}'$. The existence of a fixed-point in \mathcal{S}' follows from the Schauder fixed point theorem. \square

B. Proofs of Optimality Theorems

We start with two preliminary lemmas.

Lemma 1. *Given a message-passing equilibrium V and an allocation decision x^* , the following three conditions are equivalent:*

(i) *For every activity a , the allocation $x_{\mathcal{R}(a)}^*$ uniquely maximizes the activity manager's problem*

$$(B.1) \quad \begin{array}{ll} \text{maximize} & U_a(x_{\mathcal{R}(a)}) \triangleq u_a(x_{\mathcal{R}(a)}) + \sum_{r \in \mathcal{R}(a)} V_{r \rightarrow a}(x_{ar}) \\ \text{subject to} & x_{ar} \in \mathcal{X}_r, \end{array} \quad \forall r \in \mathcal{R}(a).$$

(ii) *For every resource r , the allocation $x_{\mathcal{A}(r)}^*$ uniquely maximizes the optimization problem*

$$(B.2) \quad \begin{array}{ll} \text{maximize} & U_r(x_{\mathcal{A}(r)}) \triangleq \sum_{a \in \mathcal{A}(r)} V_{a \rightarrow r}(x_{ar}) \\ \text{subject to} & \sum_{a' \in \mathcal{A}(r)} x_{a'r} \leq b_r, \\ & x_{a'r} \in \mathcal{X}_r, \end{array} \quad \forall a' \in \mathcal{A}(r).$$

(iii) *For every activity a and every resource $r \in \mathcal{R}(a)$, the quantity x_{ar}^* uniquely maximizes the optimization problem*

$$(B.3) \quad \begin{array}{ll} \text{maximize} & U_{ar}(x_{ar}) \triangleq V_{a \rightarrow r}(x_{ar}) + V_{r \rightarrow a}(x_{ar}) \\ \text{subject to} & x_{ar} \in \mathcal{X}_r. \end{array}$$

Proof. Given an activity a and a resource $r \in \mathcal{R}(a)$, define

$$\mathcal{C}_{a \rightarrow r} \triangleq \{x_{\mathcal{R}(a) \setminus r} : x_{ar'} \in \mathcal{X}_r, \forall r' \in \mathcal{R}(a) \setminus r\}.$$

This is the set of consumption decisions of activity a for all resources except r . Given a resource r and an activity $a \in \mathcal{A}(r)$, define $\mathcal{C}_{r \rightarrow a}(x_{ar}) \triangleq \{x_{\mathcal{A}(r) \setminus a} : \sum_{a' \in \mathcal{A}(r) \setminus a} x_{a'r} \leq b_r - x_{ar}, x_{a'r} \in \mathcal{X}_r, \forall a' \in \mathcal{A}(r) \setminus a\}$. This is the set of set of feasible allocations of resource r for all activities except a , given the allocation x_{ar} to activity a . Finally, for each resource r , define $\mathcal{C}_r(x_{ar}) \triangleq \{x_{\mathcal{A}(r) \setminus a} : \sum_{a' \in \mathcal{A}(r) \setminus a} x_{a'r} \leq b_r - x_{ar}, x_{a'r} \in \mathcal{X}_r, \forall a' \in \mathcal{A}(r) \setminus a\}$.

Then, from the equilibrium equation $HV = V$, we have for every x_{ar} ,

$$(B.4) \quad \begin{aligned} \max_{x_{\mathcal{R}(a) \setminus r} \in \mathcal{C}_{a \rightarrow r}} U_a(x_{\mathcal{R}(a)}) &= U_{ar}(x_{ar}) + (FV)_{a \rightarrow r}(0), \\ \max_{x_{\mathcal{A}(r) \setminus a} \in \mathcal{C}_{r \rightarrow a}(x_{ar})} U_r(x_{\mathcal{A}(r)}) &= U_{ar}(x_{ar}) + (FV)_{r \rightarrow a}(0). \end{aligned}$$

Assume that (iii) holds. Then, each $U_{ar}(\cdot)$ is maximized uniquely by x_{ar}^* . Consider an alternative feasible allocation x' with $x'_{ar} \neq x_{ar}^*$, for some activity a and resource $r \in \mathcal{R}(a)$. By (B.4), $x'_{\mathcal{R}(a)}$ cannot maximize $U_a(\cdot)$ and $x'_{\mathcal{A}(r)}$ cannot maximize $U_r(\cdot)$, respectively. Hence, (iii) implies (i) and (ii). The rest of the implications are shown similarly. \square

Lemma 2. *Consider a message-passing equilibrium $HV = V$, where each activity manager's problem (B.1) has a unique solution, and denote the resulting allocation by x^* . Then, for each activity a and resource $r \in \mathcal{R}(a)$, this allocation maximizes the optimization problems*

$$(B.5a) \quad \begin{aligned} \text{maximize} \quad & T_{r \rightarrow a}(x_{\mathcal{A}(r)}) \triangleq \sum_{a' \in \mathcal{A}(r) \setminus a} V_{a' \rightarrow r}(x_{a'r}) - V_{r \rightarrow a}(x_{ar}) \\ \text{subject to} \quad & \sum_{a' \in \mathcal{A}(r)} x_{a'r} \leq b_r, \\ & x_{a'r} \in \mathcal{X}_r, \quad \forall a' \in \mathcal{A}(r), \end{aligned}$$

$$(B.5b) \quad \begin{aligned} \text{maximize} \quad & T_{a \rightarrow r}(x_{\mathcal{R}(a)}) \triangleq u_a(x_{\mathcal{R}(a)}) + \sum_{r' \in \mathcal{R}(a) \setminus r} V_{r' \rightarrow a}(x_{ar'}) - V_{a \rightarrow r}(x_{ar}) \\ \text{subject to} \quad & x_{ar'} \in \mathcal{X}_{r'}, \quad \forall r' \in \mathcal{R}(a). \end{aligned}$$

Proof. Note that $T_{r \rightarrow a}(x_{\mathcal{A}(r)}) = U_r(x_{\mathcal{A}(r)}) - U_{ar}(x_{ar})$ and $T_{a \rightarrow r}(x_{\mathcal{R}(a)}) = U_a(x_{\mathcal{R}(a)}) - U_{ar}(x_{ar})$. The result then follows from (B.4) and Lemma 1. \square

Consider a message-passing equilibrium V , assume that each activity manager's problem (B.1) has a unique solution, and define x^* to be the resulting allocation. Consider an alternative feasible allocation $x \in \mathcal{X}$. These allocations differ according to the set of transfers $\Delta(x, x^*)$. We can define sets $\tilde{\mathcal{A}}$ and $\tilde{\mathcal{R}}$ of, respectively, activities and resources affected by the transfers by $\tilde{\mathcal{A}} = \{a \in \mathcal{A} : \exists r \in \mathcal{R} \text{ with } x_{ar} \neq x_{ar}^*\}$ and $\tilde{\mathcal{R}} = \{r \in \mathcal{R} : \exists a \in \mathcal{A} \text{ with } x_{ar} \neq x_{ar}^*\}$. Note that we have suppressed the dependence of the sets $\tilde{\mathcal{A}}$ and $\tilde{\mathcal{R}}$ on

x and x^* for notational simplicity. We have the following theorem, from which Theorem 2 follows as an immediate corollary.

Theorem 6. *Define an undirected bipartite graph with vertices $\tilde{\mathcal{A}}$ and $\tilde{\mathcal{R}}$, and with edges according to the set of transfers $\Delta(x, x^*)$. Then:*

- (i) *If the bipartite graph contains at most one cycle per connected component, then $U(x^*) \geq U(x)$.*
- (ii) *If, in addition, the graph contains a connected component that does not have a cycle, $U(x^*) > U(x)$.*

Proof. Recall the objective functions $U_a(\cdot)$, $U_r(\cdot)$, and $U_{ar}(\cdot)$ defined by the equilibrium V through the optimization problems (B.1), (B.2), and (B.3), respectively. The system objective $U(\cdot)$ can be written as

$$U(x) = \sum_{a \in \mathcal{A}} U_a(x_{\mathcal{R}(a)}) + \sum_{r \in \mathcal{R}} U_r(x_{\mathcal{A}(r)}) - \sum_{a \in \mathcal{A}} \sum_{r \in \mathcal{R}(a)} U_{ar}(x_{ar}).$$

We have the decomposition

$$\begin{aligned} U(x^*) - U(x) &= \sum_{a \in \tilde{\mathcal{A}}} [U_a(x_{\mathcal{R}(a)}^*) - U_a(x_{pa})] + \sum_{r \in \tilde{\mathcal{R}}} [U_r(x_{\mathcal{A}(r)}^*) - U_r(x_{\mathcal{A}(r)})] \\ &\quad - \sum_{(a,r) \in \Delta(x, x^*)} [U_{ar}(x_{ar}^*) - U_{ar}(x_{ar})]. \end{aligned}$$

By the hypothesis of the theorem, we can associate each edge $(a, r) \in \Delta(x, x^*)$ in the bipartite graph with either the vertex $a \in \tilde{\mathcal{A}}$ or the vertex $r \in \tilde{\mathcal{R}}$, in a way such that each vertex is associated with at most a single edge. Then,

$$\begin{aligned} U(x^*) - U(x) &= \sum_{a \in \tilde{\mathcal{A}}_1} [U_a(x_{\mathcal{R}(a)}^*) - U_{a\sigma(a)}(x_{a\sigma(a)}^*) - (U_a(x_{\mathcal{R}(a)}) - U_{a\sigma(a)}(x_{a\sigma(a)}))] \\ &\quad + \sum_{r \in \tilde{\mathcal{R}}_1} [U_r(x_{\mathcal{A}(r)}^*) - U_{\tau(r)r}(x_{\tau(r)r}^*) - (U_r(x_{\mathcal{A}(r)}) - U_{\tau(r)r}(x_{\tau(r)r}))] \\ &\quad + \sum_{a \in \tilde{\mathcal{A}} \setminus \tilde{\mathcal{A}}_1} [U_a(x_{\mathcal{R}(a)}^*) - U_a(x_{\mathcal{R}(a)})] + \sum_{r \in \tilde{\mathcal{R}} \setminus \tilde{\mathcal{R}}_1} [U_r(x_{\mathcal{A}(r)}^*) - U_r(x_{\mathcal{A}(r)})], \end{aligned}$$

where $\tilde{\mathcal{A}}_1 \subset \tilde{\mathcal{A}}$ and $\tilde{\mathcal{R}}_1 \subset \tilde{\mathcal{R}}$ are sets of vertices which have been associated with edges, and the maps $\sigma: \tilde{\mathcal{A}}_1 \rightarrow \tilde{\mathcal{R}}$ and $\tau: \tilde{\mathcal{R}}_1 \rightarrow \tilde{\mathcal{A}}$ define the associations. Observe that, by the unique optimality assumption and Lemmas 1 and 2, $U_r(x_{\mathcal{A}(r)}^*) > U_r(x_{\mathcal{A}(r)})$, $U_a(x_{\mathcal{R}(a)}^*) > U_a(x_{\mathcal{R}(a)})$, $U_r(x_j^*) - U_{ar}(x_{ar}^*) \geq U_r(x_j) - U_{ar}(x_{ar})$, and $U_a(x_{\mathcal{R}(a)}^*) - U_{ar}(x_{ar}^*) \geq U_a(x_{\mathcal{R}(a)}) - U_{ar}(x_{ar})$. Thus $U(x^*) \geq U(x)$. Under the additional assumption of Part (ii), the sets $\tilde{\mathcal{A}} \setminus \tilde{\mathcal{A}}_1$ and $\tilde{\mathcal{R}} \setminus \tilde{\mathcal{R}}_1$ cannot both be empty. Hence, $U(x^*) > U(x)$. \square

Theorem 4. Consider a message-passing equilibrium with concave and Lipschitz continuous messages. The resulting allocation of resources is globally optimal for the system manager's problem.

Proof. Consider a message-passing equilibrium V with concave and Lipschitz continuous messages, and let x^* be the associated allocation. Assume that x^* lies in the interior of the domain of $U(\cdot)$. By (Rockafellar, 1970, Theorem 27.4), for each resource r and activity a , there must exist a supergradient $d^{ar} \in \partial u_a(x_{\mathcal{R}(a)}^*)$ so that we have the first order conditions for the optimization problem (B.5b),

(B.6a)

$$(B.6b) \quad \begin{aligned} d_{ar}^{ar} - \frac{d^+}{dx_{ar}} V_{a \rightarrow r}(x_{ar}^*) &\leq 0, & d_{ar}^{ar} - \frac{d^-}{dx_{ar}} V_{a \rightarrow r}(x_{ar}^*) &\geq 0, \\ d_{ar'}^{ar} + \frac{d^+}{dx_{ar'}} V_{r' \rightarrow a}(x_{ar'}^*) &\leq 0, \quad \forall r' \in \mathcal{R}(a) \setminus r, & d_{ar'}^{ar} - \frac{d^-}{dx_{ar'}} V_{r' \rightarrow a}(x_{ar'}^*) &\geq 0, \quad \forall r' \in \mathcal{R}(a) \setminus r. \end{aligned}$$

Similarly, let $\lambda_{ar}^* \geq 0$ be a shadow price to the optimization problem (B.5a). Then,

(B.7a)

$$-\frac{d^+}{dx_{ar}} V_{r \rightarrow a}(x_{ar}^*) - \lambda_{ar}^* \leq 0, \quad -\frac{d^-}{dx_{ar}} V_{r \rightarrow a}(x_{ar}^*) - \lambda_{ar}^* \geq 0,$$

(B.7b)

$$\frac{d^+}{dx_{a'r}} V_{a' \rightarrow r}(x_{a'r}^*) - \lambda_{ar}^* \leq 0, \quad \forall a' \in \mathcal{A}(r) \setminus a, \quad \frac{d^-}{dx_{a'r}} V_{a' \rightarrow r}(x_{a'r}^*) - \lambda_{ar}^* \geq 0, \quad \forall a' \in \mathcal{A}(r) \setminus a.$$

Then, by (B.6a) and (B.7a),

$$\frac{d^-}{dx_{ar}} V_{a \rightarrow r}(x_{ar}^*) \leq d_{ar}^{ar} \leq \frac{d^+}{dx_{ar}} V_{a \rightarrow r}(x_{ar}^*), \quad \frac{d^-}{dx_{ar}} V_{r \rightarrow a}(x_{ar}^*) \leq -\lambda_{ar}^* \leq -\frac{d^+}{dx_{ar}} V_{r \rightarrow a}(x_{ar}^*).$$

By concavity of $V_{a \rightarrow r}(\cdot)$ and $V_{r \rightarrow a}(\cdot)$,

$$(B.8) \quad \frac{d}{dx_{ar}} V_{a \rightarrow r}(x_{ar}^*) = d_{ar}^{ar}, \quad \frac{d}{dx_{ar}} V_{r \rightarrow a}(x_{ar}^*) = -\lambda_{ar}^*,$$

where the derivatives must exist since the directional derivatives are equal. By (B.7b), and (B.8), we have $\lambda_{ar}^* = d_{a'r}^{a'r}$, for all $a' \in \mathcal{A}(r) \setminus a$. Then, must have $\lambda_{ar}^* = p_r^*$, for some vector $p^* \in \mathbb{R}_+^{\mathcal{R}}$, and, using (B.6b), also $d_{ar'}^{ar} = p_{r'}^*$, for all $r' \in \mathcal{R}(a) \setminus r$.

Define the vector dU by $(dU)_{ar} = p_r^*$, for each $a \in \mathcal{A}$ and $r \in \mathcal{R}(a)$. Then, $dU \in \partial U(x^*)$ is a supergradient of $U(\cdot)$ at x^* , the vector p^* is a shadow price vector for the system manager's optimization problem, and the allocation x^* is globally optimal. The case where x^* is on the boundary of the domain of $U(\cdot)$ is handled similarly. \square

Theorem 5. *Let x^* be the globally optimal allocation for the system manager's problem and let p^* be a supporting price vector. Suppose that $U(\cdot)$ is differentiable at x^* . Consider a message-passing equilibrium V with concave and Lipschitz continuous messages. Then, for each activity a and resource r ,*

$$\frac{d}{dx_{ar}}V_{a \rightarrow r}(x_{ar}^*) = p_r^*, \quad \frac{d}{dx_{ar}}V_{r \rightarrow a}(x_{ar}^*) = -p_r^*,$$

where the existence of the above derivatives is guaranteed. Thus,

$$\frac{\partial}{\partial x_{ar}}u_a(x_{\mathcal{R}(a)}^*) = \frac{d}{dx_{ar}}V_{a \rightarrow r}(x_{ar}^*) = -\frac{d}{dx_{ar}}V_{r \rightarrow a}(x_{ar}^*) = p_r^*.$$

Proof. This follows by the same argument as in Theorem 4, and the fact that if $U(\cdot)$ is differentiable at x^* , $\partial U(x^*) = \{\nabla U(x^*)\}$. □

References

Rockafellar, R. T. 1970. *Convex Analysis*. Princeton University Press, Princeton, NJ.